Machine Learning I

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Machine Learning for Computer Vision TU Dresden



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Contents.

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- ► This problem is introduced as an unsupervised learning problem w.r.t. constrained data.

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Definition. A strict order on A is a binary relation $< \subseteq A \times A$ that satisfies the following conditions:

$$\forall a \in A: \quad \neg a < a \tag{1}$$

$$\forall \{a, b\} \in \binom{A}{2} \colon \quad a < b \text{ xor } b < a$$
 (2)

$$\forall \{a, b, c\} \in \binom{A}{3}: \quad a < b \land b < c \Rightarrow a < c \tag{3}$$

Lemma. The strict orders on A are characterized by the bijections $\alpha:\{0,\dots,|A|-1\}\to A.$ For any such bijection, consider the order $<_{\alpha}$ such that

$$\forall a, b \in A: \quad a < b \iff \alpha^{-1}(a) < \alpha^{-1}(b) . \tag{4}$$

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$$\forall a, b \in A: \quad a < b \Leftrightarrow \alpha^{-1}(a) < \alpha^{-1}(b) . \tag{4}$$

Lemma. The strict orders on A are characterized by those

$$y: \{(a,b) \in A \times A \mid a \neq b\} \to \{0,1\}$$
 (5)

that satisfy the following conditions:

$$\forall a \in A \ \forall b \in A \setminus \{a\} \colon \quad y_{ab} + y_{ba} = 1 \tag{6}$$

$$\forall a \in A \ \forall b \in A \setminus \{a\} \ \forall c \in A \setminus \{a,b\} \colon \quad y_{ab} + y_{bc} - 1 \le y_{ac} \tag{7}$$

Constrained Data

We reduce the problem of learning and inferring orders to the problem of learning and inferring decisions, by defining constrained data (S,X,x,Y) with

$$S = \{(a,b) \in A \times A \mid a \neq b\}$$

$$\mathcal{Y} = \left\{ y \in \{0,1\}^S \mid \forall a \in A \ \forall b \in A \setminus \{a\} : \quad y_{ab} + y_{ba} = 1 \right\}$$

$$\forall a \in A \ \forall b \in A \setminus \{a\} \ \forall c \in A \setminus \{a,b\} :$$

$$y_{ab} + y_{bc} - 1 \leq y_{ac} \right\}$$

$$(8)$$

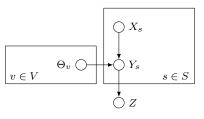
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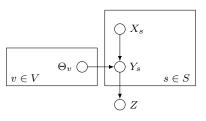
- \blacktriangleright We consider a finite, non-empty set V , called a set of attributes, and the attribute space $X=\mathbb{R}^V$
- ▶ We consider **linear functions**. Specifically, we consider $\Theta = \mathbb{R}^V$ and $f:\Theta \to \mathbb{R}^X$ such that

$$\forall \theta \in \Theta \ \forall \hat{x} \in \mathbb{R}^V : \quad f_{\theta}(\hat{x}) = \sum_{v \in V} \theta_v \, \hat{x}_v = \langle \theta, \hat{x} \rangle .$$
 (10)



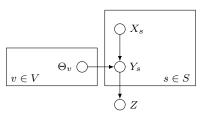
Random Variables

▶ For any $(a,b) = s \in S = E$, let X_s be a random variable whose value is a vector $x_s \in \mathbb{R}^V$, the **attribute vector** of s.



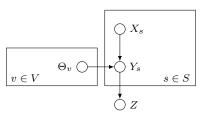
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- For any $(a,b)=s\in S$, let Y_s be a random variable whose value is a binary number $y_s\in\{0,1\}$, called the **decision** placing a before b.



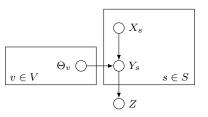
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- For any $v \in V$, let Θ_v be a random variable whose value is a real number $\theta_v \in \mathbb{R}$, a parameter of the function we seek to learn.



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- For any $v \in V$, let Θ_v be a random variable whose value is a real number $\theta_v \in \mathbb{R}$, a parameter of the function we seek to learn.
- ▶ Let Z be a random variable whose value is a subset $\mathcal{Z} \subseteq \{0,1\}^S$ called the set of **feasible decisions**. For ordering, we are interested in $\mathcal{Z} = \mathcal{Y}$, the set of characteristic functions of strict orders on A.



Factorization

$$P(X, Y, Z, \Theta) = P(Z \mid Y) \prod_{s \in S} P(Y_s \mid X_s, \Theta) \prod_{v \in V} P(\Theta_v) \prod_{s \in S} P(X_s)$$

Factorization

► Supervised learning:

$$P(\Theta \mid X, Y, Z)$$

Factorization

Supervised learning:

$$\begin{split} P(\Theta \mid X,Y,Z) &= \frac{P(X,Y,Z,\Theta)}{P(X,Y,Z)} \\ &= \frac{P(Z \mid Y) \, P(Y \mid X,\Theta) \, P(X) \, P(\Theta)}{P(Z \mid X,Y) \, P(X,Y)} \\ &= \frac{P(Z \mid Y) \, P(Y \mid X,\Theta) \, P(X) \, P(\Theta)}{P(Z \mid Y) \, P(X,Y)} \\ &= \frac{P(Y \mid X,\Theta) \, P(X) \, P(\Theta)}{P(X,Y)} \\ &\propto P(Y \mid X,\Theta) \, P(\Theta) \\ &= \prod_{s \in S} P(Y_s \mid X_s,\Theta) \prod_{v \in V} P(\Theta_v) \end{split}$$

Factorization

► Inference:

$$P(Y \mid X, Z, \theta)$$

Factorization

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$$P(Y \mid X, Z, \theta) = \frac{P(X, Y, Z, \Theta)}{P(X, Z, \Theta)}$$

$$= \frac{P(Z \mid Y) P(Y \mid X, \Theta) P(X) P(\Theta)}{P(X, Z, \Theta)}$$

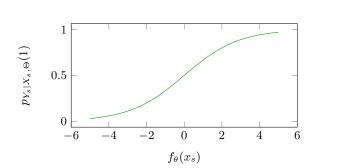
$$\propto P(Z \mid Y) P(Y \mid X, \Theta)$$

$$= P(Z \mid Y) \prod_{s \in S} P(Y_s \mid X_s, \Theta)$$

Distributions

► Sigmoid distribution

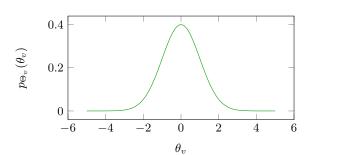
$$\forall s \in S: \qquad p_{Y_s|X_s,\Theta}(1) = \frac{1}{1 + 2^{-f_{\theta}(x_s)}}$$
 (11)



Distributions

▶ Normal distribution with $\sigma \in \mathbb{R}^+$:

$$\forall v \in V: \qquad p_{\Theta_v}(\theta_v) = \frac{1}{\sigma\sqrt{2\pi}} e^{-\theta_v^2/2\sigma^2} \tag{12}$$



Distributions

► Uniform distribution on a subset

$$\forall \mathcal{Z} \subseteq \{0,1\}^S \ \forall y \in \{0,1\}^S \quad p_{Z|Y}(\mathcal{Z},y) \propto \begin{cases} 1 & \text{if } y \in \mathcal{Z} \\ 0 & \text{otherwise} \end{cases}$$

Note that $p_{Z|Y}(\mathcal{Y},y)$ is non-zero iff the labeling $y\colon S\to\{0,1\}$ defines an order on A.

Lemma. Estimating maximally probable parameters θ , given attributes x and decisions y, i.e.,

$$\underset{\theta \in \mathbb{R}^{V}}{\operatorname{argmax}} \quad p_{\Theta|X,Y,Z}(\theta,x,y,\mathcal{Y})$$

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Proof. Analogous to the case of deciding, we obtain:

$$\underset{\theta \in \mathbb{R}^{V}}{\operatorname{argmax}} \quad p_{\Theta|X,Y,Z}(\theta, x, y, \mathcal{Y})$$

$$= \underset{\theta \in \mathbb{R}^{V}}{\operatorname{argmin}} \quad \sum_{s \in S} \left(-y_{s} f_{\theta}(x_{s}) + \log \left(1 + 2^{f_{\theta}(x_{s})} \right) \right) + \frac{\log e}{2\sigma^{2}} \|\theta\|_{2}^{2}.$$

Lemma. Estimating maximally probable decisions y, given attributes x, given the set of feasible decisions \mathcal{Y} , and given parameters θ , i.e.,

$$\underset{y \in \{0,1\}^S}{\operatorname{argmax}} \quad p_{Y|X,Z,\Theta}(y,x,\mathcal{Y},\theta) \tag{13}$$

assumes the form of the linear ordering problem:

$$\underset{y \colon S \to \{0,1\}}{\operatorname{argmin}} \quad \sum_{s \in S} (-\langle \theta, x_s \rangle) \, y_s \tag{14}$$

subject to
$$\forall a \in A \ \forall b \in A \setminus \{a\}: \quad y_{ab} + y_{ba} = 1$$
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Theorem. The linear ordering problem is NP-hard.

The linear ordering problem has been studied intensively. A comprehensive survey is by Martí and Reinelt (2011). Pioneering research is by Grötschel (1984).

We define two **local search algorithms** for the linear ordering problem.

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For simplicity, we define $c:S \to \mathbb{R}$ such that

$$\forall s \in S \colon \quad c_s = -\langle \theta, x_s \rangle \tag{17}$$

and write the (linear) cost function $\varphi:\{0,1\}^S \to \mathbb{R}$ such that

$$\forall y \in \{0,1\}^S \colon \quad \varphi(y) = \sum_{s \in S} c_s \, y_s \tag{18}$$

Greedy transposition algorithm:

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Definition. For any bijection $\alpha:\{0,\ldots,|A|-1\}\to A$ and any $j,k\in\{0,\ldots,|A|-1\}$, let $\mathrm{transpose}_{jk}[\alpha]$ the bijection obtained from α by swapping α_j and α_k , i.e.

$$\forall l \in \{0, \dots, |A| - 1\} \colon \text{ transpose}_{jk}[\alpha](l) = \begin{cases} \alpha_k & \text{if } l = j \\ \alpha_j & \text{if } l = k \\ \alpha_l & \text{otherwise} \end{cases}$$
 (19)

$$\begin{split} \alpha' &= \mathsf{greedy\text{-}transposition}(\alpha) \\ \mathsf{choose}\ (j,k) &\in \underset{0 \leq j' < k' < |A|}{\operatorname{argmin}} \varphi(y^{\mathsf{transpose}_{j'k'}[\alpha]}) - \varphi(y^{\alpha}) \\ \mathsf{if}\ \varphi(y^{\mathsf{transpose}_{jk}[\alpha]}) - \varphi(y^{\alpha}) < 0 \\ \alpha' &:= \mathsf{greedy\text{-}transposition}(\mathsf{transpose}_{jk}[\alpha]) \\ \mathsf{else} \\ \alpha' &:= \alpha \end{split}$$

Greedy transposition using the technique of Kernighan and Lin (1970)

```
\alpha' = \mathsf{greedy\text{-}transposition\text{-}kl}(\alpha)
\alpha^0 := \alpha
\delta_0 := 0
J_0 := \{0, \ldots, |A| - 1\}
                                                                                                                                          (build sequence of swaps)
repeat
      \mathsf{choose}\ (j,k) \in \mathrm{argmin}\ \varphi(y^{\mathsf{transpose}}{j'k'}^{[\alpha^t]}) - \varphi(u^{\alpha^t})
                         \{(i',k')\in J_{+}^{2}|i'< k'\}
      \begin{split} \alpha^{t+1} &:= \mathsf{transpose}_{jk}[\alpha_t] \\ \delta_{t+1} &:= \varphi(\boldsymbol{y}^{\alpha^{t+1}}) - \varphi(\boldsymbol{y}^{\alpha^t}) < 0 \end{split}
       J_{t+1} := J_t \setminus \{j, k\}
                                                                                                                                  (move \alpha_i and \alpha_k only once)
until |J_t| < 2
\hat{t} := \min \underset{t' \in \{0, \dots, |A|\}}{\operatorname{argmin}} \sum_{\tau=0}^{t'} \delta_{\tau}
                                                                                                                                                (choose sub-sequence)
if \sum_{\tau=0}^{\iota} \delta_{\tau} < 0
      \alpha' := \text{greedy-transposition-kl}(\alpha^{\hat{t}})
                                                                                                                                                                         (recurse)
else
      \alpha' := \alpha
                                                                                                                                                                     (terminate)
```

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- ► The inference problem assumes the form of the NP-hard linear ordering problem
- ► Local search algorithms for tackling this problem are greedy transposition and greedy transposition using the technique of Kernighan and Lin.