

Machine Learning I

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Binary Decision Trees

Contents. This part of the course is about the supervised learning of binary decision trees.

- ▶ We introduce the problem as a specialization of supervised learning by defining labeled data, a family of functions, a regularizer and a loss function.
- ▶ We prove that the problem is NP-hard, by relating it to the exact cover by 3-sets problem.

Binary Decision Trees

We consider labeled data with **binary features**. More specifically, we consider some finite, non-empty set V , called the set of features, and labeled data $T = (S, X, x, y)$ such that $X = \{0, 1\}^V$. Hence:

$$x: S \rightarrow \{0, 1\}^V$$

$$y: S \rightarrow \{0, 1\}$$

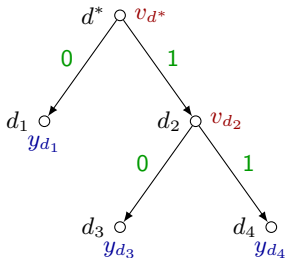
Example.

7	1	0	6	2	0	6	2
3	8	0	0	8	8	4	7
5	8	7	3	9	0	5	8
1	5	0	2	8	4	2	3
0	4	3	9	8	2	1	8
5	0	1	6	6	5	5	2
1	7	7	1	2	3	7	3
6	3	7	6	6	1	4	0

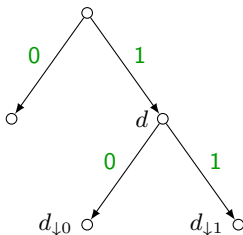
Binary Decision Trees

Definition. A tuple $(V, Y, D, D', d^*, E, \delta, v, y)$ is called a V -variate Y -valued **binary decision tree** (BDT) if and only if the following conditions hold:

1. $V \neq \emptyset$ is finite (called the set of **variables**)
2. $Y \neq \emptyset$ is finite (called the set of **values**)
3. $(D \cup D', E)$ is a finite, non-empty directed binary tree with root d^*
4. every $d \in D'$ is a leaf
5. $\delta: E \rightarrow \{0, 1\}$
6. every $d \in D$ has precisely two out-edges, $e = (d, d')$, $e' = (d, d'')$, such that $\delta(e) = 0$ and $\delta(e') = 1$
7. $v: D \rightarrow V$
8. $y: D' \rightarrow Y$



Definition. For any BDT $(V, Y, D, D', d^*, E, \delta, v, y)$, any $d \in D$ and any $j \in \{0, 1\}$, let $d_{\downarrow j} \in D \cup D'$ the unique node such that $e = (d, d_{\downarrow j}) \in E$ and $\delta(e) = j$.



Definition. For any BDT $\theta = (V, Y, D, D', d^*, E, \delta, v, y)$ and any $d \in D \cup D'$, the tuple $\theta[d] = (V, Y, D_2, D'_2, d, E', \delta', v', y')$ is called the **binary decision subtree** of θ rooted at d iff

- ▶ $(D_2 \cup D'_2, E')$ is the subtree of $(D \cup D', E)$ rooted at d
- ▶ δ', v' and y' are the restrictions of δ, v and y to the subsets D_2, D'_2 and E'

Lemma. For any BDT $\theta = (V, Y, D, D', d^*, E, \delta, v, y)$ and any $d \in D \cup D'$, the binary decision subtree $\theta[d]$ is itself a V -variate Y -valued BDT.

Definition. For any BDT $\theta = (V, Y, D, D', d^*, E, \delta, v, y)$, the function defined by θ is the $f_\theta : \{0, 1\}^V \rightarrow Y$ such that $\forall x \in \{0, 1\}^V$:

$$\begin{aligned} f_\theta(x) &= \begin{cases} y(d^*) & \text{if } D = \emptyset \\ f_{\theta[d^*_{\downarrow 0}]}(x) & \text{if } D \neq \emptyset \wedge x_{v(d^*)} = 0 \\ f_{\theta[d^*_{\downarrow 1}]}(x) & \text{if } D \neq \emptyset \wedge x_{v(d^*)} = 1 \end{cases} \\ &= \begin{cases} y(d^*) & \text{if } D = \emptyset \\ (1 - x_{v(d^*)})f_{\theta[d^*_{\downarrow 0}]}(x) + x_{v(d^*)}f_{\theta[d^*_{\downarrow 1}]}(x) & \text{otherwise} \end{cases} \end{aligned}$$

Remark. The set Θ of V -variate $Y = \{0, 1\}$ -valued BDTs can be identified with a subset of V -variate disjunctive normal forms.

Definition. For any BDT $\theta = (V, Y, D, D', d^*, E, \delta, v, y)$, the **depth** of θ is the $R(\theta) \in \mathbb{N}$ such that

$$R(\theta) = \begin{cases} 0 & \text{if } D = \emptyset \\ 1 + \max\{R(\theta[d_{\downarrow 0}^*]), R(\theta[d_{\downarrow 1}^*])\} & \text{otherwise} \end{cases} . \quad (1)$$

Binary Decision Trees

Definition. For any labeled data $T = (S, X, x, y)$ with $X = \{0, 1\}^V$, the set Θ of all V -variate $\{0, 1\}$ -valued BDTs, the family $f : \Theta \rightarrow \{0, 1\}^X$ of functions defined by these BDTs, the depth R of BDTs as a regularizer, the 0-1-loss L and any $\lambda \in \mathbb{R}_0^+$:

- The instance of the **supervised learning** problem of BDTs has the form

$$\min_{\theta \in \Theta} \quad \lambda R(\theta) + \sum_{s \in S} L(f_\theta(x_s), y_s) \quad (2)$$

- The **separation problem** of BDTs has the form

$$\inf_{\theta \in \Theta} \quad R(\theta) \quad (3)$$

$$\text{subject to} \quad \forall s \in S : \quad f_\theta(x_s) = y_s \quad (4)$$

- For any $m \in \mathbb{N}$, the **separability problem** of BDTs is to decide whether there exists a BDT $\theta \in \Theta$ such that

$$R(\theta) \leq m \quad (5)$$

$$\forall s \in S : \quad f_\theta(x_s) = y_s \quad (6)$$

Remark. separability $\leq_p$ separation $\leq_p$ supervised learning¹.

¹ $\leq_p$: Karp reduction.

Next, we show that *separability* is NP-hard by reducing the **exact cover by 3-sets** problem, using a construction by Haussler (1988).

Definition. Let S be any set and $\Sigma \subseteq 2^S$. Σ is called a **cover** of S if and only if

$$\bigcup_{\sigma \in \Sigma} \sigma = S . \quad (7)$$

A cover of S is called **exact** if and only if

$$\forall \{\sigma, \sigma'\} \in \binom{\Sigma}{2}: \quad \sigma \cap \sigma' = \emptyset . \quad (8)$$

Definition. Let S be any set and $\Sigma \subseteq 2^S$. Deciding whether there exists a $\Sigma' \subseteq \Sigma$ such that Σ' is an exact cover of S is called the instance of the **exact cover problem** w.r.t. S and Σ . If, in addition, $|S|$ is an integer multiple of 3 and any $U \in \Sigma$ is such that $|U| = 3$, the instance of the exact cover problem wrt. S and Σ is also called the **exact cover by 3-sets problem** wrt. S and Σ .

Proof. For any instance (S', Σ) of the exact cover by 3-sets problem and the $n \in \mathbb{N}$ such that $|S'| = 3n$, we construct the instance of the separability problem of BDTs such that

- ▶ $V = \Sigma$
- ▶ $S = S' \cup \{0\}$
- ▶ $x : S \rightarrow \{0, 1\}^\Sigma$ such that $x_0 = 0$ and

$$\forall s \in S' \forall \sigma \in \Sigma: \quad x_s(\sigma) = \begin{cases} 1 & \text{if } s \in \sigma \\ 0 & \text{otherwise} \end{cases} \quad (9)$$

- ▶ $y : S \rightarrow \{0, 1\}$ such that $y_0 = 0$ and $\forall s \in S': y_s = 1$.
- ▶ $m = n$

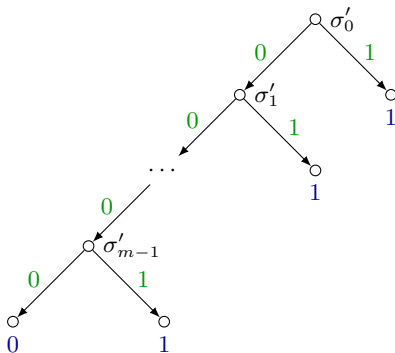
We show that the instance the exact cover problem has a solution iff the instance of the separability problem of BDTs has a solution.

Binary Decision Trees

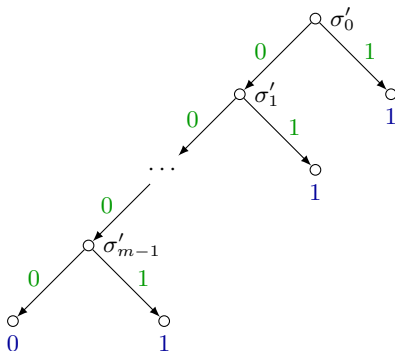
($\Rightarrow$) Let $\Sigma' \subseteq \Sigma$ a solution to the instance of the exact cover problem.

Consider any order on Σ' and the bijection $\sigma' : n \rightarrow \Sigma'$ induced by this order.

We show that the BDT θ depicted below solves the instance of the separability problem of BDTs.



Binary Decision Trees



The BDT satisfies $R(\theta) = m$.

The BDT decides the labeled data correctly because

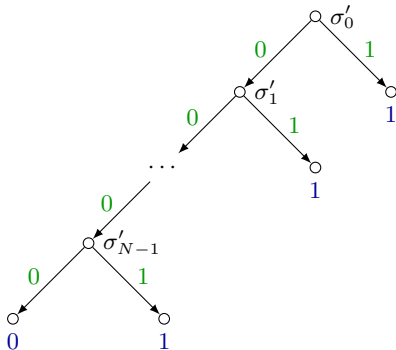
- $f_\theta(x_0) = 0 = y_0$
- At each of the m interior nodes, three additional elements of S' are mapped to 1. Thus, all $3m$ many elements $s \in S'$ are mapped to 1. That is $\forall s \in S': f_\theta(x_s) = 1 = y_s$.

Binary Decision Trees

($\Leftarrow$) Let $\theta = (V, Y, D, D', d^*, E, \delta, \sigma, y')$ a BDT that solves the instance of the separability problem of BDTs.

W.l.o.g., we assume, for any interior node $d \in D$, that $d_{\downarrow 1}$ is a leaf and $y'(d_{\downarrow 1}) = 1$.

Hence, θ is of the form depicted below.



Binary Decision Trees

Therefore:

$$\forall x \in X: \quad f_{\theta}(x) = \begin{cases} 1 & \text{if } \exists j \in N: x(\sigma_j) = 1 \\ 0 & \text{otherwise} \end{cases} \quad (10)$$

Thus,

$$\forall s \in S: \quad f_{\theta}(x_s) = \begin{cases} 1 & \text{if } \exists j \in N: s \in \sigma_j \\ 0 & \text{otherwise} \end{cases}, \quad (11)$$

by definition of x in (9).

Consequently,

$$\bigcup_{j=0}^{N-1} \sigma_j = S', \quad (12)$$

by definition of y such that $\forall s \in S': y_s = 1$.

Binary Decision Trees

Moreover, $N = m$, because

$$3m = |S'| \stackrel{(12)}{=} \left| \bigcup_{j=0}^{N-1} \sigma_j \right| \leq \sum_{j=0}^{N-1} |\sigma_j| = \sum_{j=0}^{N-1} 3 = 3N \stackrel{(5)}{\leq} 3m .$$

Therefore:

$$\forall k \in N \ \forall l \in N \setminus \{k\}: \quad \sigma_k \cap \sigma_l = \emptyset \quad (13)$$

Thus,

$$\bigcup_{j=0}^{N-1} \sigma_j$$

is a solution to the instance of the exact cover by 3-sets problem defined by (S', Σ) , by (12) and (13).

□

Summary:

- ▶ Supervised learning of BDTs is hard.
- ▶ More specifically, the NP-hard exact cover by 3-sets problem is reducible to the separability problem of BDTs, by construction of Haussler data.

Topics of upcoming exercises:

- ▶ A heuristic algorithm for the supervised learning of BDTs
- ▶ Supervised learning of disjunctive normal forms